

GRID DIAGRAMS AND THE OZSVÁTH–SZABÓ TAU-INVARIANT

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ABSTRACT. We use grid diagrams to investigate the Ozsváth–Szabó concordance invariant τ , and to prove that $|\tau(K_1) - \tau(K_2)| \leq g$, whenever there is a genus g knot cobordism joining K_1 to K_2 . This leads to an entirely grid diagram-based proof of Kronheimer–Mrowka’s theorem, formerly known as the Milnor conjecture.

1. Introduction

Links inside S^3 can be represented by a combinatorial structure called grid diagrams, and these grid diagrams can then be used to study various properties of the links. Grid diagrams first appeared as arc-presentations in [1], and are also equivalent to the square-bridge positions of [2], the Legendrian realizations of [3], the asterisk presentations of [4] and the fences of [5]. They have been used to classify essential tori in the complement of non-split links [6], to define certain Legendrian and transverse knot invariants [7], and to describe an algorithm to detect the unknot [8]. Many properties of grid diagrams have been studied in great detail in [9].

Quite recently, it was observed in [10] that grid diagrams can also be used to study a family of knot invariants and link invariants called knot Floer homology, originally defined for knots in [11, 12], and extended for links in [13]. Knot Floer homology is a powerful knot invariant, which generalizes the Alexander polynomial and can detect the knot genus [14], and fiberedness [15]. However, it was originally defined using holomorphic geometry, and it is an interesting endeavor to find combinatorial reinterpretations of various aspects of the theory using grid diagrams.

The aspect of knot Floer homology that we will study here is the Ozsváth–Szabó knot invariant τ , as defined in [12, 16]. It was shown in [16] that the absolute value of τ is a lower bound for the four-ball genus, and it can be used to prove a theorem due to Kronheimer and Mrowka [17], formerly known as the Milnor conjecture, that the unknotting number of the torus knot $T(p, q)$ is $\frac{(p-1)(q-1)}{2}$. In this paper, we will study the definition of τ in terms of grid diagrams, we will compute τ for torus knots using grid diagrams, and we will give a grid diagram-based proof of the fact that $|\tau|$ is a lower bound for the four-ball genus which similar in spirit to Rasmussen’s proof for the s invariant [18], thereby giving a new combinatorial proof of the Milnor’s conjecture.

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2. Knot cobordisms

Throughout this paper, the terms *knots* and *links* will mean oriented knots and oriented links inside S^3 . A *link cobordism* from a link L_1 to a link L_2 is a properly embedded oriented surface S inside $S^3 \times [0, 1]$, such that $\partial S \cap (S^3 \times \{0\}) = -L_1 \times \{0\}$ and $\partial S \cap (S^3 \times \{1\}) = L_2 \times \{1\}$. A link cobordism joining a knot to another knot is called a *knot cobordism*. After a small isotopy of the cobordism S inside $S^3 \times [0, 1]$ relative to the boundary, we can assume that the second projection $p_2: S^3 \times [0, 1] \rightarrow \mathbb{R}$, restricted to S , is a Morse function. We will call this Morse function $p_2|_S$, the *time function* t . The index 0, index 1 and index 2 critical points of the time function are called *births*, *saddles* and *deaths*. In a saddle, either two link components *merge* to form a single link component, or a link component *splits* to form two link components. A link cobordism S joining L_1 to L_2 is called a *concordance* if S is homeomorphic to $L_1 \times [0, 1]$. A concordance where the time function does not have any critical points is called an *isotopy*. Two links are said to be *isoptic* if there is an isotopy joining them. Two links are said to be *concordant* if there is a concordance joining them.

There is a concordance invariant τ for knots [16], such that if there is a connected genus g cobordism from K_1 to K_2 , then $|\tau(K_1) - \tau(K_2)| \leq g$. The *four-ball genus* of a knot K is the smallest integer $g^*(K)$ such that there is a connected genus $g^*(K)$ cobordism from K to the unknot. The *unknotting number* $u(K)$ of a knot K is the smallest number of crossing changes that needs to be done to convert it to the unknot. However, a crossing change is a particular type of a connected genus 1 cobordism. Therefore, for any knot K , we have $u(K) \geq g^*(K) \geq |\tau(K)|$.

In the subsequent sections, we will start with the definition of τ using grid diagrams, and directly prove that the inequality $|\tau(K_1) - \tau(K_2)| \leq g$ holds whenever there is a connected genus g cobordism joining K_1 to K_2 . Representing torus knots by grid diagrams, we will compute τ and produce an explicit unknotting sequence, and thereby give a new and completely grid diagram-based proof of Milnor's conjecture $u(T(p, q)) = \frac{(p-1)(q-1)}{2}$. As a preparatory move, let us prove the following lemma about knot cobordisms.

Lemma 2.1. *If S is a connected knot cobordism from K_1 to K_2 , then S can be isotoped to a cobordism S' inside $S^3 \times [0, 1]$ relative to the boundary, preserving the number of births, saddles and deaths throughout the isotopy, such that for S' , all the births happen at time $\frac{1}{4}$, all the saddles happen at time $\frac{1}{2}$, and all the deaths happen at time $\frac{3}{4}$. Furthermore, we can ensure that S' restricted to $S^3 \times [0, \frac{1}{4})$ and S' restricted to $S^3 \times (\frac{3}{4}, 1]$ are both product cobordisms.*

Proof. We would like to think of cobordisms as movies with time running from 0 to 1. The still at time t is $S \cap (S^3 \times \{t\})$; therefore, all but finitely many of the stills are links in S^3 . We start with some link, and as the movie plays, for most of the time, we simply isotope the link. However, at finitely many points in time, we can have births, saddles or deaths, as shown in Figure 1.

Clearly, we can isotope S , while preserving the number of births, saddles and deaths, to ensure that all the deaths happen at the very end. Just before some death is about to happen, intervene, and keep the relevant unknot component alive. Since the unknot, thus kept alive, is an unknot supported inside a very small ball, it behaves like a point, and therefore generically does not interfere with the rest of the cobordism.

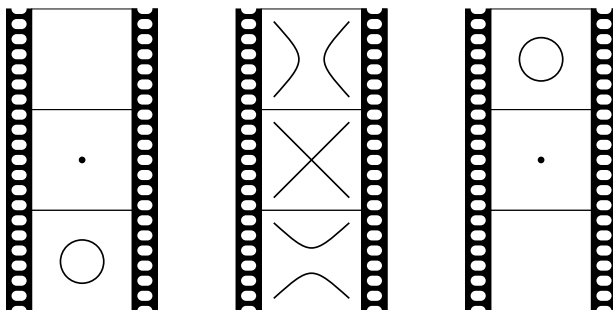


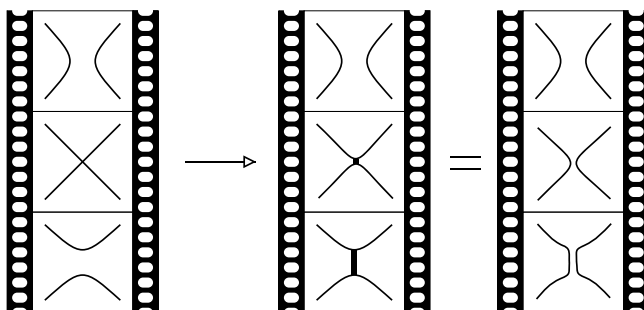
FIGURE 1. Birth, saddle, and death.

Therefore, we can postpone all the deaths, until all that is left of the cobordism is a product cobordism, and then have all the deaths. Similarly, we can ensure that all the births happen at the very beginning. Thus we can assume that the cobordism S restricted to either $S^3 \times [0, \frac{1}{4})$ or $S^3 \times (\frac{3}{4}, 1]$ is a product cobordism, and all the births happen at time $\frac{1}{4}$, and all the deaths happen at time $\frac{3}{4}$.

Now we want to isotope the cobordism inside $S^3 \times [\frac{1}{4}, \frac{3}{4}]$, relative to the boundary, so as to ensure that all the saddles happen at the same time. By reparametrizing time if necessary, assume that all the saddles happen before time $t = \frac{1}{2}$. We will now describe how to postpone all the saddles until that point, and then make all the saddles happen simultaneously.

During the movie for S , if at some point in time a saddle happens, then in the movie for the new cobordism S' , immediately after that point, attach an untwisted band to the two strands near the saddle and add an 1-handle to the link along that band. Therefore, in the modified cobordism S' , the saddle has not yet taken place. However, while watching the movie for S' , if we ignore all these new bands and the associated 1-handles, it will look exactly like the movie for S . This modification from S to S' is shown in Figure 2, with the bands being denoted by thick lines.

In the movie for S' , move the endpoints of the bands as prescribed by the movie for S , and move the bands themselves in any fashion, while ensuring that they stay disjoint from each other and from the rest of the link. Then at time $t = \frac{1}{2}$, after

FIGURE 2. Modifying the cobordism S to get the cobordism S' by delaying a saddle.

we have encountered all the saddles of S , and after we have attached bands for each one of them in S' , actually do all the saddles for S' . The saddles have the effect of cutting open all the bands, which can then deformation retract to their endpoints on the link. After that, the movie for S' agrees the movie for S . \square

3. Grid diagrams

The best reference for this section is [19]. Many of the definitions and theorems that we are about to mention in this section, are treated in great detail in that paper. However, for the sake of completeness, let us still go through some of the basic definitions and state some of the basic properties of grid diagrams.

3.1. Grid diagrams for S^3 . An *index- n S^3 -grid diagram* $\mathfrak{G} = (T, \alpha, \beta, O)$ is a picture of the following type on a torus T : α and β are two n -component embedded multicurves on T ; each α -circle is transverse to each β -circle, and they intersect each other at one point; the n components of $T \setminus \alpha$ are called the *horizontal annuli*; the n components of $T \setminus \beta$ are called the *vertical annuli*; O is a formal sum of n markings on T , such that each horizontal annulus contains one O -marking and each vertical annulus contains one O -marking; for some $0 \leq k \leq n$, exactly $(n - k)$ of the O -markings are designated *special* and are often denoted by \emptyset ; the other O -markings are called *normal* O -markings are numbered O_1, \dots, O_k .

A *generator* x is a formal sum of n points on T , often called *x -coordinates*, such that each α -circle contains one x -coordinate and each β -circle contains one x -coordinate. The set of all the $n!$ generators is denoted by $\mathcal{G}_{\mathfrak{G}}$. Given two generators $x, y \in \mathcal{G}_{\mathfrak{G}}$ which differ in exactly two coordinates, a *rectangle* joining x to y is an embedded rectangle $R \subset T$ such that: the sides of R lie on $\alpha \cup \beta$; the top-right and bottom-left corners of R are x -coordinates and the top-left and bottom-right corners of R are y -coordinates, or in other words, $\partial(\partial R|_{\alpha}) = y - x$; R does not contain any other x -coordinates; and R does not contain any special O -marking. For $x, y \in \mathcal{G}_{\mathfrak{G}}$, the set $\mathcal{R}_{\mathfrak{G}}(x, y)$ is defined to be empty if x and y do not differ in exactly two coordinates, or else, it is defined to be the set of all rectangles joining x to y . Given a rectangle $R \in \mathcal{R}_{\mathfrak{G}}(x, y)$, the number $n_i(R)$ is defined to be 1 if R contains O_i , and is defined to be 0 otherwise.

To each generator we can associate an integer-valued grading M called the *Maslov grading* in the following way: the torus is cut up along some α -circle and some β -circle and identified with $[0, n) \times [0, n)$, such that the α -circles become the lines $[0, n) \times \{i\}$ for $i \in \{0, \dots, n-1\}$, and the β -circles become the lines $\{i\} \times [0, n)$ for $i \in \{0, \dots, n-1\}$; let \mathcal{J} be the bilinear form on the singular 0-chains of \mathbb{R}^2 , such that, if $p = (p_1, p_2)$ and $q = (q_1, q_2)$ are two points in \mathbb{R}^2 , $\mathcal{J}(p, q) = \frac{1}{2}$ if $(p_1 - q_1)(p_2 - q_2) > 0$ and $\mathcal{J}(p, q) = 0$ otherwise; for any generator $x \in \mathcal{G}_{\mathfrak{G}}$, the Maslov grading is defined as $M(x) = \mathcal{J}(x - O, x - O) + 1$.

The *grid chain complex* $CF_{\mathfrak{G}}$ over \mathbb{F}_2 is defined in the following way: it is freely generated over $\mathbb{F}_2[U_1, \dots, U_k]$ by the elements of $\mathcal{G}_{\mathfrak{G}}$; the Maslov grading is extended to $CF_{\mathfrak{G}}$ by declaring the Maslov grading of each U_i to be (-2) ; the homological grading is simply the Maslov grading; the boundary map ∂ is U_i -equivariant and for

any $x \in \mathcal{G}_{\mathfrak{G}}$, it is given by

$$\partial x = \sum_{y \in \mathcal{G}_{\mathfrak{G}}} y \sum_{R \in \mathcal{R}_{\mathfrak{G}}(x,y)} \prod_i U_i^{n_i(R)}.$$

Theorem 3.1. [10, 19] *If k , the number of normal O -markings, is less than n , then the homology of $CF_{\mathfrak{G}}$ is isomorphic, as graded $\mathbb{F}_2[U_1, \dots, U_k]$ -modules, to $\otimes^{n-k-1}(\mathbb{F}_2 \oplus \mathbb{F}_2[-1])$, where each U_i acts trivially on the right hand side, each \mathbb{F}_2 lives in grading zero and $[i]$ denotes a grading shift by i .*

In light of the above theorem, the number $k + 1 - n$ is often called the *smallest Maslov grading* since it is the smallest grading in which the homology of $CF_{\mathfrak{G}}$ is supported; furthermore, the rank of the homology in the smallest Maslov grading is always one.

3.2. Grid diagrams for links. An *index- n link-grid diagram* $\mathfrak{L} = (T, \alpha, \beta, O, X)$ is a picture on a torus T such that: X is a formal sum of n points on the torus; if $\mathfrak{f}(\mathfrak{L}) = (T, \alpha, \beta, O)$ is the diagram obtained from \mathfrak{L} by forgetting the X -markings, then $\mathfrak{f}(\mathfrak{L})$ is an index- n S^3 -grid diagram; furthermore, each horizontal annulus contains some X -marking and each vertical annulus contains some X -marking.

Given an index- n link-grid diagram \mathfrak{L} , we can produce n^2 links in the following way: cut the torus T along some α -circle and some β -circle to identify it with $[0, n) \times [0, n)$; in every horizontal strip, join the X -marking to the O -marking by a horizontal line segment, and in every vertical strip join the O -marking to the X -marking by a vertical line segment, with the understanding that if there is a square containing both an X -marking and an O -marking, then we put a small unknot in that square; and finally at every crossing, declare the vertical segment to be the overpass. These n^2 links, thus obtained, are all isotopic to one another; therefore, a link-grid diagram represents a link isotopy class. Whenever we say that a link L is represented by a link-grid diagram \mathfrak{L} , we mean that \mathfrak{L} represents the link isotopy class that contains L . Call \mathfrak{L} *tight*, if every link component in \mathfrak{L} contains exactly one special O -marking.

Lemma 3.1. [9] *Every link can be represented by link-grid diagrams.*

In a link-grid diagram, generators can be endowed with a $\frac{1}{2}\mathbb{Z}$ -valued grading A called the *Alexander grading* as follows: the torus is cut up along an α -circle and β -circle such that it can be identified with $[0, n) \times [0, n)$; if \mathcal{J} is the bilinear form on the 0-chains of \mathbb{R}^2 from before, then for any generator $x \in \mathcal{G}_{\mathfrak{f}(\mathfrak{L})}$, the Alexander grading is defined as

$$\begin{aligned} A(x) &= \mathcal{J}(x - \frac{1}{2}(X + O), X - O) - \frac{n-1}{2} \\ &= \mathcal{J}(x, X) - \mathcal{J}(x, O) - \frac{1}{2}\mathcal{J}(X, X) + \frac{1}{2}\mathcal{J}(O, O) - \frac{n-1}{2}. \end{aligned}$$

This can be extended to an Alexander grading on $CF_{\mathfrak{f}(\mathfrak{L})}$ by declaring the Alexander grading of each U_i to be (-1) . An astute reader will observe that our definition of Alexander grading differs from the usual definition of Alexander grading [19] by an additive constant of $\frac{l-1}{2}$, where l is the number of link components; therefore, the two definitions agree for knots. The boundary map ∂ does not increase this Alexander grading. This leads to the following definition of the *Alexander filtration* on the grid

chain complex: for every $a \in \frac{1}{2}\mathbb{Z}$, the filtration level $\mathcal{F}_{\mathfrak{L}}(a) \subseteq CF_{\mathfrak{f}(\mathfrak{L})}$ is defined to be the subcomplex supported in Alexander grading a or less.

Define $\tau_{\mathfrak{L}}$ to be smallest $a \in \frac{1}{2}\mathbb{Z}$ such that the map induced on homology from the inclusion $\mathcal{F}_{\mathfrak{L}}(a) \hookrightarrow CF_{\mathfrak{f}(\mathfrak{L})}$ is non-trivial.

Theorem 3.2. [10] *If a knot K is represented by a tight link-grid diagram \mathfrak{L} , then $\tau(K) = \tau_{\mathfrak{L}}$.*

This is actually very close to the original definition of τ . Combining this fact with the main result from [16], we get the following:

Theorem 3.3. [16, 10] *For tight link-grid diagrams \mathfrak{L} that represent knots, $\tau_{\mathfrak{L}}$ depends only on the isotopy class of the knot. If we define $\tau(K)$ to be equal to $\tau_{\mathfrak{L}}$ for any tight link-grid diagram representing K , then $|\tau(K_1) - \tau(K_2)| \leq g$, whenever there is a connected genus g knot cobordism from K_1 to K_2 .*

The proof of this theorem requires the holomorphic techniques of [16]. We will bypass those methods, and give a new proof of this theorem using only grid diagrams. That is one of our main results.

3.3. Moves on S^3 -grid diagrams. In this subsection, we will describe certain S^3 -grid moves which convert an S^3 -grid diagram \mathfrak{G}_1 to another S^3 -grid diagrams \mathfrak{G}_2 , and in each case, we will define chain maps from $CF_{\mathfrak{G}_1}$ to $CF_{\mathfrak{G}_2}$. Given a link cobordism from a link L represented by a link-grid diagram \mathfrak{L} to a link L' represented by a link-grid diagram \mathfrak{L}' , we will be able to construct a sequence of link-grid diagrams $\mathfrak{L} = \mathfrak{L}_0, \mathfrak{L}_1, \dots, \mathfrak{L}_{m-1}, \mathfrak{L}_m = \mathfrak{L}'$, such that for each i , $\mathfrak{f}(\mathfrak{L}_i)$ and $\mathfrak{f}(\mathfrak{L}_{i+1})$ will be related by one of the following S^3 -grid moves. Therefore, we will have chain maps $CF_{\mathfrak{f}(\mathfrak{L}_i)} \rightarrow CF_{\mathfrak{f}(\mathfrak{L}_{i+1})}$, and by composing, we will get a chain map from $CF_{\mathfrak{f}(\mathfrak{L})}$ to $CF_{\mathfrak{f}(\mathfrak{L}')}$, which we will use to find a relation between $\tau_{\mathfrak{L}}$ and $\tau_{\mathfrak{L}'}$.

S^3 -grid move (1). $\mathfrak{G}_1 = \mathfrak{G}_2$. In this case, we define the chain map to be identity, which clearly preserves Maslov grading and is a quasi-isomorphism.

S^3 -grid move (2). \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by renumbering the normal O -markings. If there are exactly k normal O -markings, which are renumbered by some permutation $\sigma \in \mathfrak{S}_k$, then the chain map sends $\prod_i U_i^{m_i} x$ to $\prod_i U_{\sigma(i)}^{m_i} x$. Although this map is not U_i -equivariant, it preserves the Maslov grading and is a quasi-isomorphism.

S^3 -grid move (3). \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by a *commutation*. There are two types of commutations: a horizontal commutation where we interchange the O -markings in two adjacent horizontal annuli, or a vertical commutation where we interchange the O -markings in two adjacent vertical annuli. In either case, we represent both \mathfrak{G}_1 and \mathfrak{G}_2 by a single diagram \mathfrak{G} on the torus T , and the chain map is defined by counting certain pentagons in \mathfrak{G} . For example, in a horizontal commutation, as illustrated in Figure 3, the shaded pentagon contributes to the chain map. As shown in [19, Section 3.1], the pentagon map also preserves the Maslov grading and is a quasi-isomorphism.

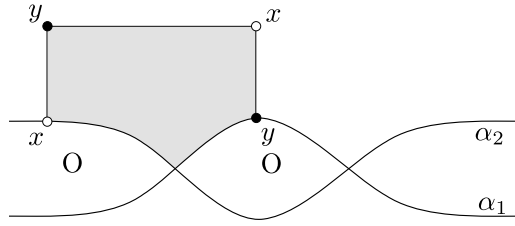


FIGURE 3. The diagram \mathfrak{G} . The grid diagram \mathfrak{G}_i is obtained from \mathfrak{G} by deleting the the circle labeled α_i . If $x \in \mathcal{G}_{\mathfrak{G}_1}$ is represented by the white circles and if $y \in \mathcal{G}_{\mathfrak{G}_2}$ is represented by the black circles, then the shaded pentagon contributes a coefficient of y for the chain map evaluated at x .

S^3 -grid move (4). \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by a *normal stabilization* or a *normal destabilization*. In a normal destabilization, we start with an index- $(n+1)$ S^3 -grid diagram \mathfrak{G}_1 with exactly $(k+1)$ normal O -markings, and we get the index- n S^3 -grid diagram \mathfrak{G}_2 by deleting O_{k+1} and then deformation retracting the closure of the horizontal annulus through O_{k+1} to an α -circle and deformation retracting the closure of the vertical annulus through O_{k+1} to a β -circle. A normal stabilization is the reverse process of a normal destabilization. Let us assume that \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by a normal destabilization; we will describe four chain maps: two chain maps d_{11} and d_{22} from $CF_{\mathfrak{G}_1}$ to $CF_{\mathfrak{G}_2}$, and two chain maps s_{11} and s_{22} from $CF_{\mathfrak{G}_2}$ to $CF_{\mathfrak{G}_1}$.

In the S^3 -grid diagram \mathfrak{G}_2 , let the α -circle that comes from deformation retracting the horizontal annulus be numbered α_0 , and let the β -circle that comes from deformation retracting the vertical annulus be numbered β_0 . In the S^3 -grid diagram \mathfrak{G}_1 , let the two α -circles on the boundary of that horizontal annulus be numbered α_1 and α_2 , such that α_1 lies just below α_2 , and let the two β -circles on the boundary of that vertical annulus be numbered β_1 and β_2 , such that β_1 lies just to left of β_2 . This is shown in the first part of Figure 4.

Let us define four injective maps $F_{ij}: \mathcal{G}_{\mathfrak{G}_2} \rightarrow \mathcal{G}_{\mathfrak{G}_1}$, for $i, j \in \{1, 2\}$. For defining F_{ij} , we identify the α -circles of \mathfrak{G}_1 , except α_i , with the α -circles of \mathfrak{G}_2 in the natural way, and we identify the β -circles of \mathfrak{G}_2 , except β_j , with the β -circles of \mathfrak{G}_2 in the natural way. Under these identifications, a generator $x \in \mathcal{G}_{\mathfrak{G}_2}$ produces a formal sum of n points in \mathfrak{G}_1 . We define $F_{ij}(x)$ to be that formal sum plus $\alpha_i \cap \beta_j$.

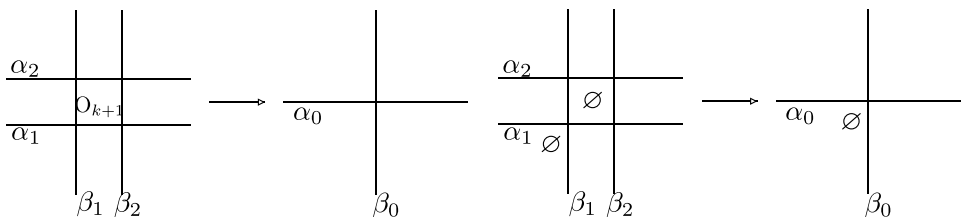


FIGURE 4. The two types of destabilizations.

The destabilization map d_{11} is precisely the snail map F^R as defined in [19, Section 3.2]. It is a homomorphism of $\mathbb{F}_2[U_1, \dots, U_k]$ -modules, and for $x \in \mathcal{G}_{\mathfrak{S}_1}$, it is defined as

$$d_{11}(U_{k+1}^m x) = U^m \sum_{y \in \mathcal{G}_{\mathfrak{S}_2}} y \sum_{D \in \mathcal{S}_1(x, F_{11}(y), \alpha_1 \cap \beta_1)} \prod_{1 \leq i \leq k} U_i^{n_i(D)},$$

where: $U = 0$ if the horizontal annulus just below α_0 contains a special O -marking, and $U = U_j$ if the horizontal annulus contains the normal O -marking O_j ; $\mathcal{S}_1(x, z, p)$ is the set of all Type (1) snail-like domains centered at p joining x to z , as illustrated in the bottom row of [19, Figure 13] or the top row of Figure 5; and $n_i(D)$ is the number of times D passes through O_i . As shown in [19], this map preserves the Maslov grading and is a quasi-isomorphism.

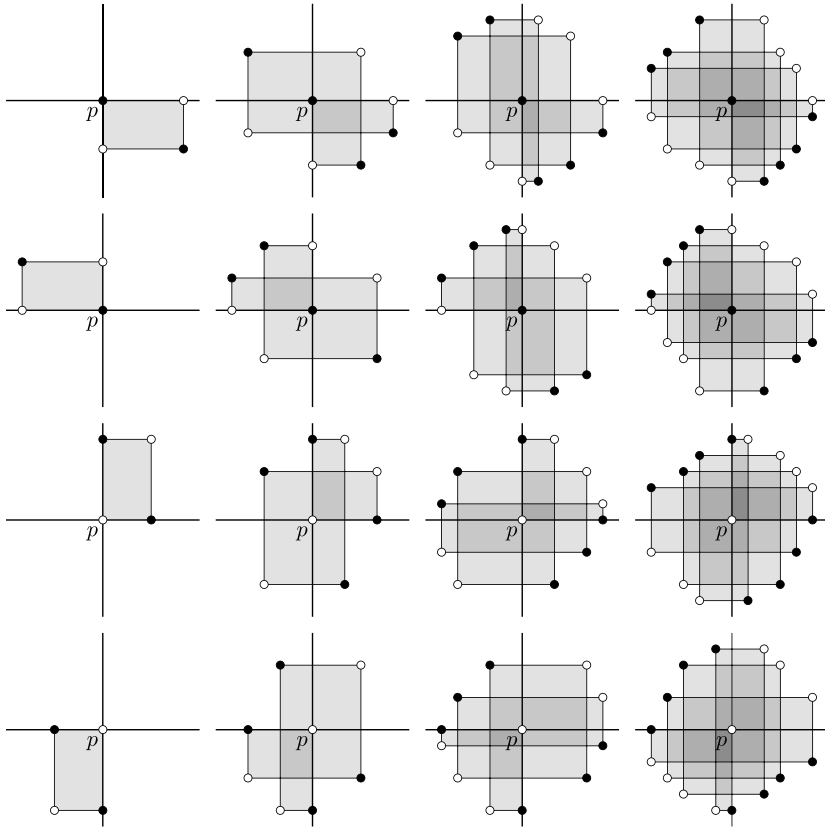


FIGURE 5. The different types of snail-like domains. The i th row displays some of the domains in $\mathcal{S}_i(x, z, p)$, where the x - and z -coordinates that are not disjoint from the domains are represented by the white and black circles, respectively. We always assume that none of the snail-like domains pass through any of the special O -markings.

The destabilization map d_{22} is obtained by rotating all the diagrams by 180° . Stated more precisely,

$$d_{22}(U_{k+1}^m x) = U^m \sum_{y \in \mathcal{G}_{\mathfrak{G}_2}} y \sum_{D \in \mathcal{S}_2(x, F_{22}(y), \alpha_2 \cap \beta_2)} \prod_{1 \leq i \leq k} U_i^{n_i(D)}$$

where: $U = 0$ if the horizontal annulus just above α_0 contains a special O -marking, and $U = U_j$ if the horizontal annulus contains the normal O -marking O_j ; $\mathcal{S}_2(x, z, p)$ is the set of all Type (2) snail-like domains centered at p joining x to z , as illustrated in the second row of Figure 5; and $n_i(D)$ is the number of times D passes through O_i .

The two stabilization maps are defined similarly. Namely, for $x \in \mathcal{G}_{\mathfrak{G}_2}$ and $j \in \{1, 2\}$,

$$s_{jj}(x) = \sum_{y \in \mathcal{G}_{\mathfrak{G}_1}} y \sum_{D \in \mathcal{S}_{j+2}(F_{jj}(x), y, \alpha_j \cap \beta_j)} \prod_{1 \leq i \leq k} U_i^{n_i(D)}.$$

The proofs of Lemma 3.5 and Proposition 3.8 of [19] go through after rotating all the diagrams by $\pm 90^\circ$. Therefore, the stabilization maps are also chain maps; furthermore, they preserve the Maslov grading and are quasi-isomorphisms.

S³-grid move (5). \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by a *special destabilization*. This is like a normal destabilization, except we assume that the O -marking that is being deleted is a special O -marking. We also assume that the square immediately to the bottom-left of this special O -marking also contains a special O -marking. The situation is illustrated in the second part of Figure 4. We will reuse the notations that we had used to describe a normal destabilization.

A *hexagon* H is an embedded hexagon in \mathfrak{G}_1 with boundary lying in $\alpha \cup \beta$, which has one 270° angle at $\alpha_1 \cap \beta_1$ and five other 90° angles, such that H contains the special O -marking that is being deleted, but not the special O -marking that lies to the bottom-left of it; H does not contain any other special O -marking, and $n_i(H)$ is defined to be 1 if H contains O_i , and is defined to be 0 otherwise. A hexagon H joins a generator $x \in \mathcal{G}_{\mathfrak{G}_1}$ to a generator $z \in \mathcal{G}_{\mathfrak{G}_1}$, if H does not contain any x -coordinate in its interior, and $\partial(\partial H|_\alpha) = z - x$; the set of all hexagons joining x to z is denoted by $\mathcal{H}(x, z)$. The chain map $f: CF_{\mathfrak{G}_1} \rightarrow CF_{\mathfrak{G}_2}$ is U_i -equivariant for all i , and for $x \in \mathcal{G}_{\mathfrak{G}_1}$, it is defined as follows:

$$f(x) = \begin{cases} F_{11}^{-1}(x), & \text{if } F_{11}^{-1}(x) \neq \emptyset, \\ \sum_{y \in \mathcal{G}_{\mathfrak{G}_2}} y \sum_{H \in \mathcal{H}(x, F_{11}(y))} \prod_i U_i^{n_i(H)}, & \text{otherwise.} \end{cases}$$

It is not hard to see that this map is a chain map, which increases the Maslov grading by 1. In fact, this map is simply the map F^L of [19, Section 3.2]. It follows from [19, Proposition 3.8] that this map is surjective at the level of homology.

S³-grid move (6). \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by the following process: we assume that \mathfrak{G}_1 has exactly $(k + 1)$ normal O -markings and a 2×2 square B which contains a special O -marking on the top-left and O_{k+1} on the bottom-right; \mathfrak{G}_2 is obtained from \mathfrak{G}_1 by deleting the two O -marking in B and adding two special O -markings, one on the bottom-left and one on the top-right of B . We define the U_i -equivariant chain map $f: CF_{\mathfrak{G}_1} \rightarrow CF_{\mathfrak{G}_2}$ as follows: set $U_{k+1} = 0$; send every generator that does not have a coordinate at the center of B to zero; and send every generator with a

coordinate at the center of B to itself. It is easy to see that this is a chain map which drops Maslov grading by 1. After passing through an extra filtration, as defined in [19, Section 3.2], we see that the induced map on the homology of the associated graded object is injective. Therefore, the map f is injective at the level of homology.

3.4. Moves on link-grid diagrams. Quite like in the previous subsection, in this subsection we will analyze certain *link-grid moves*, which convert a link-grid diagram \mathfrak{L}_1 to another link-grid diagram \mathfrak{L}_2 . In all the link-grid moves that we will analyze, the two S^3 -grid diagrams $\mathfrak{f}(\mathfrak{L}_1)$ and $\mathfrak{f}(\mathfrak{L}_2)$ will already be related by one of the six S^3 -grid moves; therefore, we already have maps $CF_{\mathfrak{f}(\mathfrak{L}_1)} \rightarrow CF_{\mathfrak{f}(\mathfrak{L}_2)}$. We will simply determine the *Alexander grading shifts* in each case. The Alexander grading shift of a map $f: CF_{\mathfrak{f}(\mathfrak{L}_1)} \rightarrow CF_{\mathfrak{f}(\mathfrak{L}_2)}$ is the smallest possible $s \in \frac{1}{2}\mathbb{Z}$, such that f shifts the Alexander grading of each element by at most s . In other words, for each $a \in \frac{1}{2}\mathbb{Z}$, there is the following commuting square.

$$\begin{array}{ccc} \mathcal{F}_{\mathfrak{L}_1}(a) & \longrightarrow & \mathcal{F}_{\mathfrak{L}_2}(a+s) \\ \downarrow & & \downarrow \\ CF_{\mathfrak{f}(\mathfrak{L}_1)} & \xrightarrow{f} & CF_{\mathfrak{f}(\mathfrak{L}_2)} \end{array}$$

Link-grid move (1). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by *renumbering* the normal O -markings. This corresponds to the S^3 -grid move (2), and the Alexander grading shift is 0.

Link-grid move (2). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by a *commutation*. Commutation comes in two flavors, horizontal commutation and vertical commutation. In a horizontal commutation, we choose two adjacent horizontal annuli such that the zero-sphere obtained by projecting the markings in one annulus to the middle α -circle is unlinked from the zero-sphere obtained by projecting the markings in the other annulus to that α -circle; and then we interchange the two horizontal annuli; cf. Figure 6. In a vertical commutation, we choose two adjacent vertical annuli such that the zero-sphere obtained by projecting the markings in one annulus to the middle β -circle is unlinked from the zero-sphere obtained by projecting the markings in the other annulus to that β -circle, and then interchange the two vertical annuli; cf. [19, Figure 5]. Commutation corresponds to the S^3 -grid move (3), and as shown in [19, Lemma 3.1], the Alexander grading shift is 0.

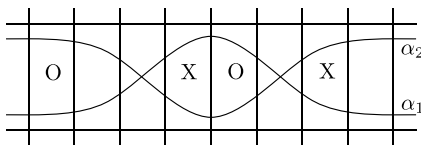


FIGURE 6. A diagram representing a horizontal commutation. The link-grid diagram \mathfrak{L}_i is obtained from this diagram by deleting the the circle labeled α_i .

Link-grid move (3). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by a *destabilization* or vice versa. In a destabilization from \mathfrak{L}_1 to \mathfrak{L}_2 , we assume that \mathfrak{L}_1 has exactly $(k+1)$ normal O -markings, and we assume that there is a 2×2 square B , three of whose squares are occupied by O_{k+1} and two X -markings. We then delete O_{k+1} and these two X -markings, and we put a new X -marking in the other square of B . We then deformation retract the horizontal annulus which contained O_{k+1} to an α -circle and deformation retract the vertical annulus which contained O_{k+1} to a β -circle to get the link-grid diagram \mathfrak{L}_2 . This move corresponds to the S^3 -grid move (4), where we use the snail maps which are centered at the center of B . As shown in [19, Lemma 3.5], the Alexander grading shift is 0.

Link-grid move (4). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by a *birth*, i.e., we assume that \mathfrak{L}_2 has exactly $(k+1)$ normal O -markings, with O_{k+1} lying in the same square as some X -marking, and \mathfrak{L}_1 is obtained from \mathfrak{L}_2 by deleting O_{k+1} and that X -marking, and then deformation retracting the horizontal and the vertical annulus through that square to an α -circle and a β -circle, respectively. This move corresponds to the S^3 -grid move (4) and it represents a birth happening in a cobordism.

We will now show that the Alexander grading shift is $(-\frac{1}{2})$. Let us reuse the notations from S^3 -grid move (4). There are two stabilization maps s_{11} and s_{22} from $CF_{\mathfrak{f}(\mathfrak{L}_1)}$ to $CF_{\mathfrak{f}(\mathfrak{L}_2)}$. We will only deal with the map s_{11} ; the map s_{22} can be dealt with in a similar fashion. Consider the map $\tilde{s}: \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_2)} \rightarrow CF_{\mathfrak{f}(\mathfrak{L}_2)}$ defined as follows:

$$\tilde{s}(x) = \sum_{y \in \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_2)}} y \sum_{D \in \mathcal{S}_{33}(x, y, \alpha_1 \cap \beta_1)} \prod_{1 \leq i \leq k} U_i^{n_i(D)}.$$

For any generator $x \in \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_1)}$, $s_{11}(x) = \tilde{s}(F_{11}(x))$. The Alexander grading shift of the map \tilde{s} is zero, since the Alexander grading shift induced by a snail-like domain D is the number of O 's minus the number of X 's (both counted with multiplicities) that are contained in D ; and O_{k+1} appears in D the same number of times as the X -marking that lies in the same square as O_{k+1} , and every other normal O -marking O_i appears with a canceling factor of U_i ; cf. [19, Proof of Lemma 3.5]. Therefore, we only have to compute the Alexander grading shift of the map $F_{11}: \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_1)} \rightarrow \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_2)}$.

Assume that the index of \mathfrak{L}_2 is $(n+1)$. Let us cut up the torus along α_1 , i.e., the α -circle that lies just below O_{k+1} , and β_2 , i.e., the β -circle that lies just to the right of O_{k+1} , to identify it with $[0, n+1) \times [0, n+1)$. The α -circles of \mathfrak{L}_2 become the horizontal lines $[0, n+1) \times \{i\}$ for $i \in \{0, \dots, n\}$ and the β -circles of \mathfrak{L}_2 become the vertical lines $\{i\} \times [0, n+1)$ for $i \in \{0, \dots, n\}$. To get from \mathfrak{L}_1 to \mathfrak{L}_2 , we start with the subsquare $[0, n) \times [1, n+1)$, we add the lowermost row $[0, n+1) \times [0, 1)$ and the rightmost column $[n, n+1) \times [0, n+1)$, and we add an X -marking and the O -marking O_{k+1} at the bottom-right square $[n, n+1) \times [0, 1)$. To get from $x \in \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_1)}$ to $F_{11}(x) \in \mathcal{G}_{\mathfrak{f}(\mathfrak{L}_2)}$, we start with a formal sum of n points in $[0, n) \times [1, n+1)$ and we add the point $(n, 0)$.

Recall that the Alexander grading of a generator y is $A(y) = \mathcal{J}(y, X) - \mathcal{J}(y, O) - \frac{1}{2}\mathcal{J}(X, X) + \frac{1}{2}\mathcal{J}(O, O) - \frac{n-1}{2}$. This process increases each of the terms $\mathcal{J}(y, X)$, $\mathcal{J}(y, O)$ and $\frac{n-1}{2}$ by $\frac{1}{2}$ and does not change the terms $\mathcal{J}(X, X)$ and $\mathcal{J}(O, O)$. Therefore, the net Alexander grading shift is $(-\frac{1}{2})$.

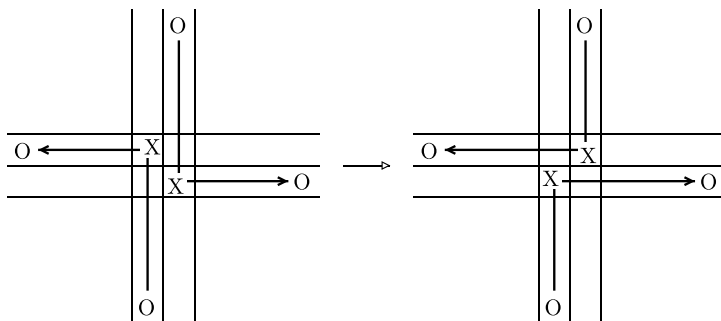


FIGURE 7. The saddle move. The two link-grid diagrams \mathfrak{L}_1 and \mathfrak{L}_2 are shown, along with the oriented links that they represent (drawn with thick lines).

Link-grid move (5). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by a *saddle*, i.e., we assume that \mathfrak{L}_1 has a 2×2 square B with two X -markings, one at the top-left and one at the bottom-right, and \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by deleting these two X -markings and putting two new ones, one at the top-right and one at the bottom-left of B . This move corresponds to the S^3 -grid move (1), and a direct computation reveals that the Alexander grading shift is $\frac{1}{2}$. This move represents a saddle happening in a cobordism, as illustrated in Figure 7.

Conversely, given a saddle from a link L_1 to a link L_2 , we can choose link-grid diagrams \mathfrak{L}_i representing L_i such that \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by a saddle move. Represent the two strands of L_1 where the saddle takes place by the configuration as shown in the first part of Figure 7; extend this to a rectilinear approximation for the rest of L_1 ; in the resulting diagram, if there is a crossing where the vertical arc is not the overpass, perform the local adjustment from [9, Figure 7] to rectify it. This produces the link-grid diagram \mathfrak{L}_1 for L_1 . Doing the saddle move to \mathfrak{L}_1 produces the link-grid diagram \mathfrak{L}_2 for L_2 .

Link-grid move (6). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by the following process: we assume that \mathfrak{L}_1 has exactly $(k+1)$ normal O -markings and a 2×2 square B with a special O -marking on the top-left and O_{k+1} on the bottom-right; \mathfrak{L}_2 obtained from \mathfrak{L}_1 by deleting these two markings and adding two new special O -markings, one at the top-right and one at the bottom-left of B . This corresponds to the S^3 -grid move (6), and the Alexander grading shift is $(-\frac{1}{2})$. This move also represents a saddle in a cobordism (it will become apparent in the proof of Theorem 3.3 why we need two types of saddle moves). Once again, it is easy to see that any saddle can be represented by such a link-grid move. Furthermore, if the saddle is a split, then \mathfrak{L}_1 is tight if and only if \mathfrak{L}_2 is tight.

Link-grid move (7). \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by a *death*, i.e., there is some special O -marking in \mathfrak{L}_1 such that the square immediately to the top-right of it contains an X -marking and a special O -marking, and \mathfrak{L}_2 is obtained from \mathfrak{L}_1 by deleting those two markings, and then deformation retracting the horizontal and vertical annulus through that square to an α -circle and a β -circle, respectively. This move corresponds

to the S^3 -grid move (5). By direct computation, we see that the Alexander grading shift is $\frac{1}{2}$. This move represents a death happening in a cobordism.

4. Main theorem

In this section, we will prove our main theorem, Theorem 3.3.

Lemma 4.1. *If \mathfrak{L}_1 and \mathfrak{L}_2 are two link-grid diagrams representing isotopic links, and if every link component in \mathfrak{L}_1 contains at most one special O -marking, and if the corresponding link component in \mathfrak{L}_2 contains the same number of special O -markings, then there is a sequence of link-grid moves of Types (1), (2) and (3), which takes \mathfrak{L}_1 to \mathfrak{L}_2 .*

Proof. This is a small extension of Cromwell's Theorem [9, Section 2]. Cromwell's theorem states that the above is true if all the O -markings are treated as equal. Therefore, we can simply take the sequence of link-grid moves as given by Cromwell, and apply them. However, we can run into the following four types of problems.

- (1) We might have to destabilize at a special O -marking.
- (2) We might have to destabilize at a normal O -marking which is not the highest numbered one.
- (3) In the final link-grid diagram and in \mathfrak{L}_2 , the special O -markings could be at different places.
- (4) The normal O -markings could be numbered differently in the final link-grid diagram and in \mathfrak{L}_2 .

The link-grid move (1), i.e., renumbering the normal O -markings, fixes two of these problems, namely the second and the fourth one. To fix the other problems, we only need a sequence of link-grid moves of Types (1), (2), and (3), which achieves the following: given a link-grid diagram \mathfrak{L} where every link component has at most one special O -marking, and given a special O -marking, we can convert that special O -marking to a normal O -marking, and convert the next O -marking in that oriented link component to a special O -marking. Assuming that there are exactly $(k - 1)$ normal O -markings in \mathfrak{L} , such a sequence of moves is shown in Figure 8. \square

Theorem 4.1. *If \mathfrak{L}_1 and \mathfrak{L}_2 are two tight link-grid diagrams representing knots K_1 and K_2 , respectively, and if there is a connected knot cobordism from K_1 to K_2 with exactly b births, s saddles and d deaths, then there is a sequence of link-grid moves taking \mathfrak{L}_1 to \mathfrak{L}_2 , such that there are exactly b link-grid moves of Type (4), $s - d$ link-grid moves of Type (5), d link-grid moves of Type (6) and d link-grid moves of Type (7), and these link-grid moves happen in this order.*

Proof. Using Lemma 2.1, we can assume that all the births happen at time $t = \frac{1}{4}$, all the saddles happen at time $t = \frac{1}{2}$, and all the deaths happen at time $t = \frac{3}{4}$. By isotoping the cobordism slightly, we can ensure that the critical points happen at distinct instants of time, and we can choose the order of the b births, the order of the s saddles and the order of the d deaths. Stated differently, given $t_1 < \cdots < t_b$ near $\frac{1}{4}$, $t_{b+1} < \cdots < t_{b+s}$ near $\frac{1}{2}$, and $t_{b+s+1} < \cdots < t_{b+s+d}$ near $\frac{3}{4}$, and given an ordering of the b births, an ordering of the s saddles and an ordering of the d deaths, we can isotope the cobordism slightly to ensure that the i th birth happens at time

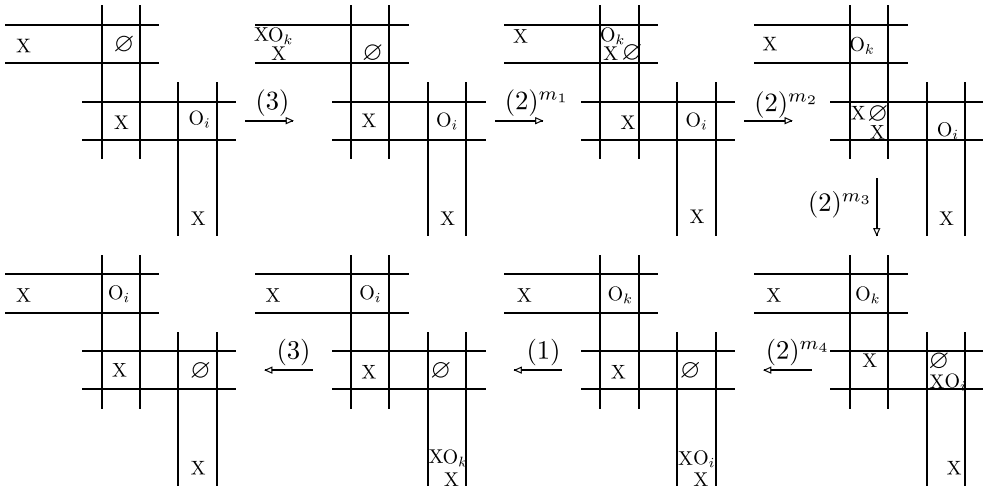


FIGURE 8. A sequence of link-grid moves.

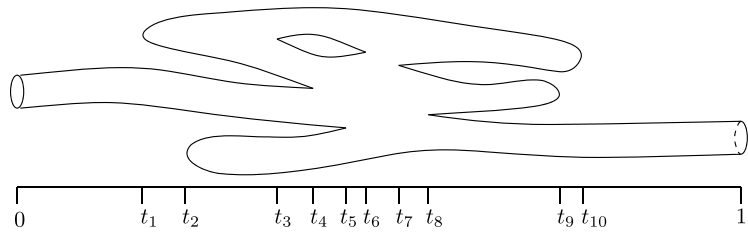


FIGURE 9. A cobordism with 2 births, 6 saddles, and 2 deaths, put in the standard form.

$t = t_i$, the i th saddle happens at time $t = t_{b+i}$ and the i th death happens at time $t = t_{b+s+i}$. Since the cobordism is connected, we order the s saddles in some way so as to guarantee that the final d saddles are all splits. In other words, we ensure that the link in the still, just after time $t = t_{b+s-d}$, is a knot. A schematic picture of a cobordism, put in this standard form, is shown in Figure 9.

Let $\mathfrak{L}_0^a = \mathfrak{L}_1$ and let $\mathfrak{L}_{b+s+d+1}^b = \mathfrak{L}_2$. For each $1 \leq i \leq b+s+d$, choose two link-grid diagrams \mathfrak{L}_i^b and \mathfrak{L}_i^a , such that: \mathfrak{L}_i^b represents the link just before time t_i ; \mathfrak{L}_i^a represents the link just after time t_i ; and \mathfrak{L}_i^a can be obtained from \mathfrak{L}_i^b by a link-grid move of Type (4), (5), (6) or (7), depending on whether $1 \leq i \leq b$, $b+1 \leq i \leq b+s-d$, $b+s-d+1 \leq i \leq b+s$ or $b+s+1 \leq i \leq b+s+d$, respectively. Observe that for each $0 \leq i \leq b+s+d$, the two link-grid diagrams \mathfrak{L}_i^a and \mathfrak{L}_{i+1}^b represent isotopic links; it is easy to see that while choosing \mathfrak{L}_i^a and \mathfrak{L}_{i+1}^b , we can ensure that the corresponding link components contain the same number of special O -markings. Therefore, by Lemma 4.1, we can convert \mathfrak{L}_i^a to \mathfrak{L}_{i+1}^b by a sequence of link-grid moves of Types (1)–(3). Putting everything together, we get the required sequence of link-grid moves that converts \mathfrak{L}_1 to \mathfrak{L}_2 . \square

Proof of Theorem 3.3. In order to prove this, we only need to show the following: if \mathfrak{L}_1 and \mathfrak{L}_2 are two tight link-grid diagrams representing knots K_1 and K_2 , respectively, and if there is a connected knot cobordism of genus g from K_1 to K_2 , then $|\tau_{\mathfrak{L}_1} - \tau_{\mathfrak{L}_2}| \leq g$.

Let us now assume that the cobordism from K_1 to K_2 has b births, d deaths and $2g + b + d$ saddles. Theorem 4.1 tells us that there is a sequence of link-grid moves of Types (1)–(7) taking \mathfrak{L}_1 to \mathfrak{L}_2 , such that there are exactly b link-grid moves of Type (4), $2g + b$ link-grid moves of Type (5), d link-grid moves of Type (6) and d link-grid moves of Type (7), and these link-grid moves happen in this order.

For link-grid moves of Types (1)–(5), the associated maps on CF preserve Maslov gradings and are quasi-isomorphisms. For link-grid move (6), the map drops Maslov grading by 1, and is injective at the level of homology. The smallest Maslov grading also drops by 1, and by Theorem 3.1, the homology of CF in the smallest Maslov grading is \mathbb{F}_2 . Therefore, the map on homology in the smallest Maslov grading is an isomorphism. Similarly, for link-grid move (7), the map increases Maslov grading by 1, and is surjective at the level of homology. However, the smallest Maslov grading also increases by 1; therefore, the map on homology in the smallest Maslov grading is also an isomorphism. Thus, the composed maps from $CF_{\mathfrak{f}(\mathfrak{L}_1)}$ to $CF_{\mathfrak{f}(\mathfrak{L}_2)}$ preserves Maslov grading and is a quasi-isomorphism in the smallest Maslov grading. However, each of the link-grid diagrams \mathfrak{L}_1 and \mathfrak{L}_2 contains exactly one special O -marking, so the smallest Maslov grading is the only Maslov grading in which the homology is supported. Therefore, the composed map is a quasi-isomorphism.

There are no Alexander grading shifts for link-grid moves (1), (2), and (3), there is an Alexander grading shift of $\frac{1}{2}$ for link-grid moves (5) and (7), and there is an Alexander grading shift of $(-\frac{1}{2})$ for link-grid moves (4) and (6). Therefore, the net Alexander grading shift is $-\frac{b}{2} + \frac{2g+b}{2} - \frac{d}{2} + \frac{d}{2} = g$. Therefore, for every $a \in \frac{1}{2}\mathbb{Z}$, we have the following commuting square.

$$\begin{array}{ccc} H_*(\mathcal{F}_{\mathfrak{L}_1}(a)) & \longrightarrow & H_*(\mathcal{F}_{\mathfrak{L}_2}(a+g)) \\ \downarrow & & \downarrow \\ H_*(CF_{\mathfrak{f}(\mathfrak{L}_1)}) = \mathbb{F}_2 & \xrightarrow{\text{Id}} & H_*(CF_{\mathfrak{f}(\mathfrak{L}_2)}) = \mathbb{F}_2 \end{array}$$

Substituting $a = \tau_{\mathfrak{L}_1}$ in the above commutative diagram, we see that the map $H_*(\mathcal{F}_{\mathfrak{L}_2}(\tau_{\mathfrak{L}_1} + g)) \rightarrow H_*(CF_{\mathfrak{f}(\mathfrak{L}_2)})$ is non-trivial; therefore, $\tau_{\mathfrak{L}_2} \leq \tau_{\mathfrak{L}_1} + g$.

However, we can view the cobordism in reverse, i.e., we can run the movie backwards, to get a connected genus g cobordism from K_2 to K_1 . That would show that $\tau_{\mathfrak{L}_1} \leq \tau_{\mathfrak{L}_2} + g$. Combining the two inequalities, we get our desired result $|\tau_{\mathfrak{L}_1} - \tau_{\mathfrak{L}_2}| \leq g$. \square

5. Applications

Given any S^3 -grid diagram \mathfrak{G} , define the *special generator* $x_{\mathfrak{G}} \in \mathcal{G}_{\mathfrak{G}}$ to be the generator whose coordinates lie at the top-left corners of the squares that contain the O -markings. It is an easy computation to show that the Maslov grading of $x_{\mathfrak{G}}$ is always zero.

Let \mathfrak{G}_n be the following index- n S^3 -grid diagram: there is exactly one special O -marking; the square containing O_{n-1} lies immediately to the top-left of the square containing the special O -marking; and for all $1 \leq i \leq n-2$, the square containing O_i lies immediately to the top-left of the square containing O_{i+1} . For this grid diagram, the special generator $x_{\mathfrak{G}_n}$ also has coordinates at the bottom-right corners of the squares that contain the O -markings.

Lemma 5.1. *The \mathbb{F}_2 -module generated by $x_{\mathfrak{G}_n}$ is a direct summand of $CF_{\mathfrak{G}_n}$.*

Proof. For any S^3 -grid diagram \mathfrak{G} , the \mathbb{F}_2 -module generated by $x_{\mathfrak{G}}$ is a quotient complex of $CF_{\mathfrak{G}}$. This is because any rectangle that joins some other generator to $x_{\mathfrak{G}}$ must pass through some O -marking. However, rectangles are not allowed to pass through the special O -markings, and whenever they pass through the normal O -markings, they pick up a U -power.

For the grid diagram \mathfrak{G}_n , we would like to show that the \mathbb{F}_2 -module generated by $x_{\mathfrak{G}_n}$ is also a subcomplex. Let $y \in \mathcal{G}_{\mathfrak{G}_n}$ be some generator that differs from $x_{\mathfrak{G}_n}$ in exactly two coordinates. There are exactly two embedded rectangles R_1 and R_2 , with boundary lying in $\alpha \cup \beta$, whose top-right and bottom-left corners are $x_{\mathfrak{G}_n}$ -coordinates, and whose top-left and bottom-right corners are y -coordinates. It is clear that none of these rectangles contain any O -markings or any $x_{\mathfrak{G}_n}$ -coordinates in their interiors. Therefore, $\partial x_{\mathfrak{G}_n} = 0$, thereby concluding the proof. \square

Lemma 5.2. *Let \mathfrak{L} be an index- n link-grid diagram that represents a knot K , such that $\mathfrak{f}(\mathfrak{L}) = \mathfrak{G}_n$. Then $\tau(K) = A(x_{\mathfrak{G}_n})$.*

Proof. The link-grid diagram \mathfrak{L} is a tight link-grid diagram representing K , therefore, $\tau(K) = \tau_{\mathfrak{L}}$, which is the smallest $a \in \frac{1}{2}\mathbb{Z}$, such that the map induced on homology from the inclusion $\mathcal{F}_{\mathfrak{L}}(a) \hookrightarrow CF_{\mathfrak{G}_n}$ is non-trivial. However, the homology of $CF_{\mathfrak{G}_n}$ is one-dimensional, carried by the direct summand which is the \mathbb{F}_2 -module generated by $x_{\mathfrak{G}_n}$. Therefore $\tau_{\mathfrak{L}} = A(x_{\mathfrak{G}_n})$. \square

Let $T(p, q)$ denote the (p, q) -torus knot. We will represent $T(p, q)$ by the following index- $(p+q)$ link-grid diagram $\mathfrak{L}_{p,q}$: $\mathfrak{f}(\mathfrak{L}_{p,q}) = \mathfrak{G}_{p+q}$; squares to the bottom-right of squares containing X -markings also contain X -markings; and the X -marking in the horizontal annulus through the special O -marking, lies p squares to the right of the special O -marking. The link-grid diagram $\mathfrak{L}_{5,3}$ is shown in Figure 10. To draw the torus knot $T(p, q)$ that $\mathfrak{L}_{p,q}$ represents or to compute Alexander gradings of specific generators, we need to identify $\mathfrak{L}_{p,q}$ with a diagram on $[0, p+q) \times [0, p+q)$. For such identifications, we always assume that the special O -marking lies in the bottom-right square.

Theorem 5.1. *There is an unknotting sequence for $T(p, q)$ with $\frac{(p-1)(q-1)}{2}$ crossing changes, and $\tau(T(p, q)) = \frac{(p-1)(q-1)}{2}$. Therefore, $u(T(p, q)) = g^*(T(p, q)) = \frac{(p-1)(q-1)}{2}$.*

Proof. Without loss of generality, let us assume that $p > q$. After identifying $\mathfrak{L}_{p,q}$ with a picture on $[0, p+q) \times [0, p+q)$, let us consider the induced knot diagram for $T(p, q)$, cf. the first part of Figure 11. In this picture, there are $\frac{q(q-1)}{2}$ crossings above the principal diagonal, and $(p-q)(q-1) + \frac{q(q-1)}{2}$ crossings below the principal diagonal,

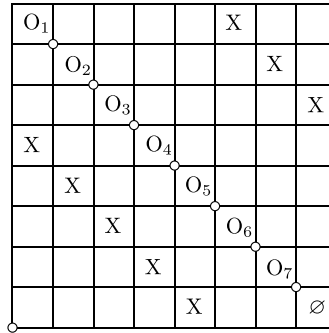


FIGURE 10. The link-grid diagram $\mathfrak{L}_{p,q}$ for $T(p,q)$, with $p = 5$ and $q = 3$. The coordinates of the special generator $x_{\mathfrak{S}_{p+q}}$ are shown.

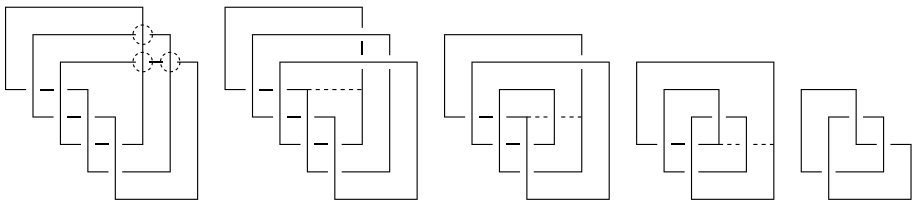


FIGURE 11. Starting at the knot diagram induced by $\mathfrak{L}_{p,q}$ with $p > q$, we do $\frac{q(q-1)}{2}$ crossing changes to get the knot diagram induced by $\mathfrak{L}_{p-q,q}$.

totalling to $p(q-1)$ crossings (thus, this is actually a minimal crossing diagram for $T(p,q)$).

Change the $\frac{q(q-1)}{2}$ crossings above the principal diagonal. Figure 11 shows how we can isotope the resulting knot diagram to get the diagram that would be induced by $\mathfrak{L}_{p-q,q}$. However, by induction, $T(p-q,q)$ can be unknotted with $\frac{(p-q-1)(q-1)}{2}$ crossing changes. Therefore, $T(p,q)$ can be unknotted with $\frac{(p-q-1)(q-1)}{2} + \frac{q(q-1)}{2} = \frac{(p-1)(q-1)}{2}$ crossing changes.

To compute $\tau(T(p,q))$, thanks to Lemma 5.2, we only need to compute the Alexander grading of the special generator $x_{\mathfrak{S}_{p+q}}$. Towards this end, let us number the coordinates of $x_{\mathfrak{S}_{p+q}}$ from left to right as x_1, \dots, x_{p+q} ; let us number the X -markings from left to right as X_1, \dots, X_{p+q} ; and let us number the special O -markings as O_{p+q} . Then,

$$\begin{aligned} A(x_{\mathfrak{S}_{p+q}}) &= \mathcal{J}(x_{\mathfrak{S}_{p+q}}, X) - \mathcal{J}(x_{\mathfrak{S}_{p+q}}, O) - \frac{1}{2}\mathcal{J}(X, X) + \frac{1}{2}\mathcal{J}(O, O) - \frac{p+q-1}{2} \\ &= \sum_{1 \leq i, j \leq p+q} \mathcal{J}(x_i, X_j) - \sum_{1 \leq i, j \leq p+q} \mathcal{J}(x_i, O_j) - \sum_{1 \leq i < j \leq p+q} \mathcal{J}(X_i, X_j) \\ &\quad + \sum_{1 \leq i < j \leq p+q} \mathcal{J}(O_i, O_j) - \frac{p+q-1}{2} \end{aligned}$$

$$\begin{aligned}
&= \sum_{j=1}^{p+q} \mathcal{J}(x_1, X_j) + \sum_{j=1}^p \sum_{i=2}^{p+q} \mathcal{J}(x_i, X_j) + \sum_{j=p+1}^{p+q} \sum_{i=2}^{p+q} \mathcal{J}(x_i, X_j) \\
&\quad - \sum_{j=1}^{p+q} \mathcal{J}(x_1, O_j) - \sum_{\substack{2 \leq i \leq p+q \\ 1 \leq j \leq p+q}} \mathcal{J}(x_i, O_j) - \sum_{\substack{1 \leq i \leq p \\ p+1 \leq j \leq p+q}} \mathcal{J}(X_i, X_j) \\
&\quad - \sum_{\substack{1 \leq i < j \leq p+q \\ i > p \text{ OR } j \leq p}} \mathcal{J}(X_i, X_j) + \sum_{1 \leq i < j \leq p+q} \mathcal{J}(O_i, O_j) - \frac{p+q-1}{2} \\
&= \frac{p+q}{2} + \frac{pq}{2} + \frac{pq}{2} - \frac{p+q}{2} - 0 - \frac{pq}{2} - 0 + 0 - \frac{p+q-1}{2} \\
&= \frac{(p-1)(q-1)}{2}.
\end{aligned}$$

Therefore, $\tau(T(p, q)) = A(x_{\mathfrak{P}_{p+q}}) = \frac{(p-1)(q-1)}{2}$. Combining our results, we get $\frac{(p-1)(q-1)}{2} \geq u(T(p, q)) \geq g^*(T(p, q)) \geq \tau(T(p, q)) = \frac{(p-1)(q-1)}{2}$, thus completing the proof. \square

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